Today: More Classical Problems

- Part 1: Leader election

- Part 2: Mutual exclusion

Termination Detection

- Detecting the end of a distributed computation
- Notation: let sender be *predecessor*, receiver be *successor*
- Two types of markers: Done and Continue
- After finishing its part of the snapshot, process Q sends a Done or a Continue to its predecessor
- Send a Done only when
  - All of Q's successors send a Done
  - Q has not received any message since it check-pointed its local state and received a marker on all incoming channels
  - Else send a Continue
- Computation has terminated if the initiator receives Done messages from everyone

Part 1: Election Algorithms

- Many distributed algorithms need one process to act as coordinator
  – Doesn't matter which process does the job, just need to pick one
- Election algorithms: technique to pick a unique coordinator (aka *leader election*)
- Examples: take over the role of a failed process, pick a master in Berkeley clock synchronization algorithm
- Types of election algorithms: Bully and Ring algorithms

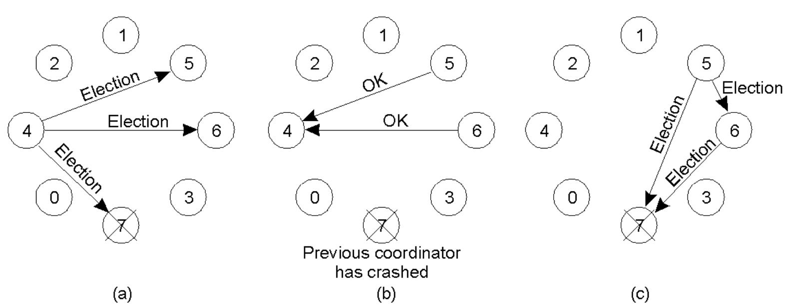
Bully Algorithm

- Each process has a unique numerical ID
- Processes know the Ids and address of every other process
- Communication is assumed reliable
- *Key Idea*: select process with highest ID
- Process initiates election if it just recovered from failure or if coordinator failed
- 3 message types: *election, OK, I won*
- Several processes can initiate an election simultaneously
  – Need consistent result
- $O(n^2)$ messages required with n processes

Bully Algorithm Details

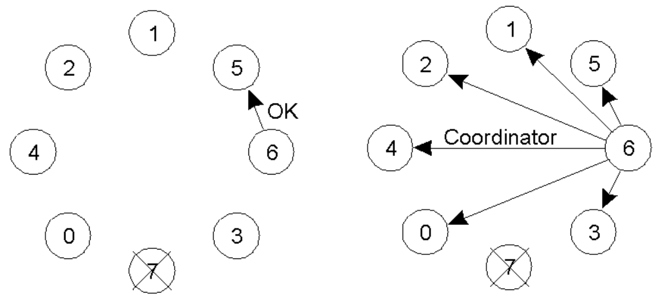
- Any process *P* can initiate an election
- *P* sends *Election* messages to all process with higher Ids and awaits *OK* messages
- If no *OK* messages, *P* becomes coordinator and sends *I won* messages to all process with lower Ids
- If it receives an *OK*, it drops out and waits for an *I won*
- If a process receives an *Election* msg, it returns an *OK* and starts an election
- If a process receives a *I won*, it treats sender an coordinator

Bully Algorithm Example



- The bully election algorithm
- Process 4 holds an election
- Process 5 and 6 respond, telling 4 to stop
- Now 5 and 6 each hold an election

Bully Algorithm Example



(d) (e)

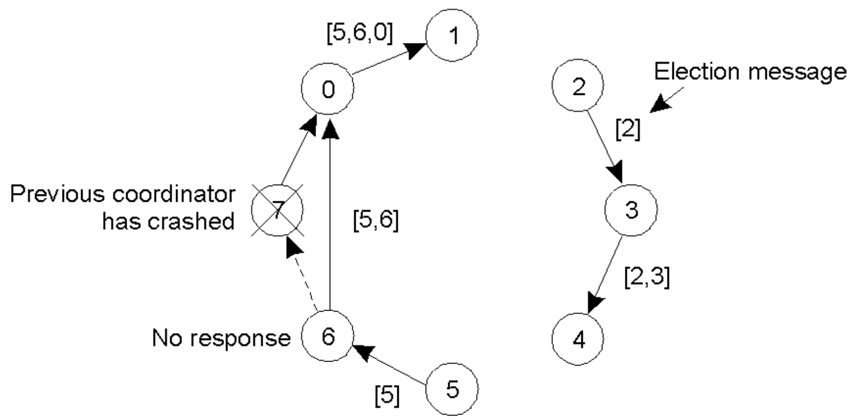
d) Process 6 tells 5 to stop

e) Process 6 wins and tells everyone

Ring-based Election

- Processes have unique Ids and arranged in a logical ring
- Each process knows its neighbors
  – Select process with highest ID
- Begin election if just recovered or coordinator has failed
- Send *Election* to closest downstream node that is alive
  – Sequentially poll each successor until a live node is found
- Each process tags its ID on the message
- Initiator picks node with highest ID and sends a coordinator message
- Multiple elections can be in progress
  – Wastes network bandwidth but does no harm

A Ring Algorithm



- Election algorithm using a ring.

Comparison

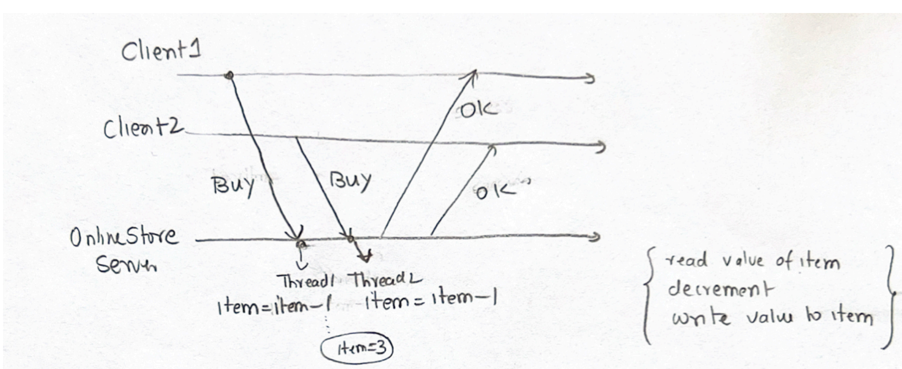
- Assume *n* processes and one election in progress
- Bully algorithm
  - Worst case: initiator is node with lowest ID
    - Triggers n-2 elections at higher ranked nodes: $O(n^2)$ msgs
  - Best case: immediate election: n-2 messages
- Ring
  - 2 (n-1) messages always

Part 2: Distributed Synchronization

- Distributed system with multiple processes may need to share data or access shared data structures
  - Use critical sections with mutual exclusion
- Single process with multiple threads
  - Semaphores, locks, monitors
- How do you do this for multiple processes in a distributed system?
  - Processes may be running on different machines
- Solution: lock mechanism for a distributed environment
  - Can be centralized or distributed

Lock Example

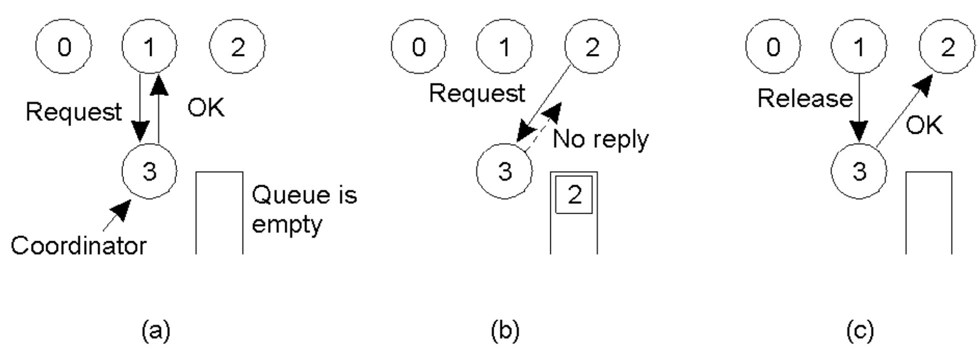
- Online store example:
  - 2 clients buy same item, need to decrement stock



Centralized Mutual Exclusion

- Assume processes are numbered
- One process is elected coordinator (highest ID process)
- Every process needs to check with coordinator before entering the critical section
- To obtain exclusive access: send request, await reply
- To release: send release message
- Coordinator:
  - Receive *request*: if available and queue empty, send grant; if not, queue request
  - Receive *release*: remove next request from queue and send grant

Mutual Exclusion: A Centralized Algorithm



a) Process 1 asks the coordinator for permission to enter a critical region. Permission is granted

b) Process 2 then asks permission to enter the same critical region. The coordinator does not reply.

c) When process 1 exits the critical region, it tells the coordinator, when then replies to 2

Properties

- Simulates centralized lock using blocking calls
- Fair: requests are granted the lock in the order they were received
- Simple: three messages per use of a critical section (request, grant, release)
- Shortcomings:
  - Single point of failure
  - How do you detect a dead coordinator?
    - A process can not distinguish between "lock in use" from a dead coordinator
      - No response from coordinator in either case
  - Performance bottleneck in large distributed systems

Decentralized Algorithm

- Use voting
- Assume n replicas and a coordinator per replica
- To acquire lock, need majority vote m > n/2 coordinators
  - Non blocking: coordinators returns OK or "no"
- Coordinator crash => forgets previous votes
  - Probability that k coordinators crash $P(k) = {}^{m}C_{k}\, p^{k}\, (1-p)^{m-k}$
  - Atleast 2m-n need to reset to violate correctness
    - $\sum_{2m-n}^{n} P(k)$

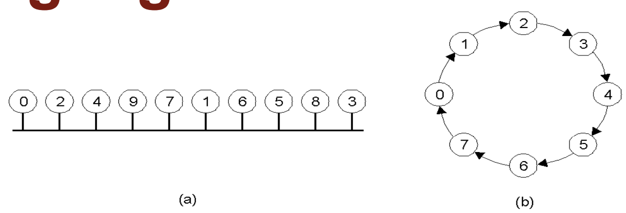
Distributed Algorithm

- [Ricart and Agrawala]: needs 2(n-1) messages
- Based on event ordering and time stamps
  - Assumes total ordering of events in the system (Lamport's clock)
- Process k enters critical section as follows
  - Generate new time stamp $TS_k = TS_k+1$
  - Send $request(k,TS_k)$ all other $n-1$ processes
  - Wait until $reply(j)$ received from all other processes
  - Enter critical section
- Upon receiving a *request* message, process j
  - Sends *reply* if no contention
  - If already in critical section, does not reply, queue request
  - If wants to enter, compare TS_j with TS_k and send reply if $TS_k<TS_j$, else queue (recall: total ordering based on multicast)

Properties

- Fully decentralized
- N points of failure!
- All processes are involved in all decisions
  - Any overloaded process can become a bottleneck

A Token Ring Algorithm



a) An unordered group of processes on a network.

b) A logical ring constructed in software.

- Use a token to arbitrate access to critical section
- Must wait for token before entering CS
- Pass the token to neighbor once done or if not interested
- Detecting token loss in non-trivial

Comparison

Algorithm	Messages per entry/ exit	Delay before entry (in message times)	Problems
Centralized	3	2	Coordinator crash
Decentralized	3mk	2m	starvation
Distributed	2 (n – 1)	2 (n – 1)	Crash of any process
Token ring	1 to ∞	0 to n – 1	Lost token, process crash

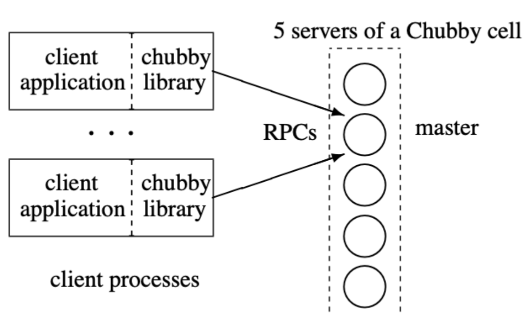
- A comparison of four mutual exclusion algorithms.

Chubby Lock Service

- Chubby: distributed lock service developed by google
  - Design for coarse-grain locking
  - uses file system abstraction for locks
  - Each Chubby cell (~5 machines) supports 10,000 servers
  - One replica is outside the data center for high availability
  - distributed file system interface for locking and sharing state
- Use cases:
  - Leader election: use locks for leader election and advertise leader
    - Grab lock, declare oneself leader
  - Coarse-grain synchronization - hold lock for hours or days

Chubby Lock Service

- Chubby cell: elect a primary
  - each replica maintains a DB
  - master initiates updates to DB
- Use file abstraction
  - file is a “named” lock
  - reader-writer locks



- Primary can fail
  - Triggers new election