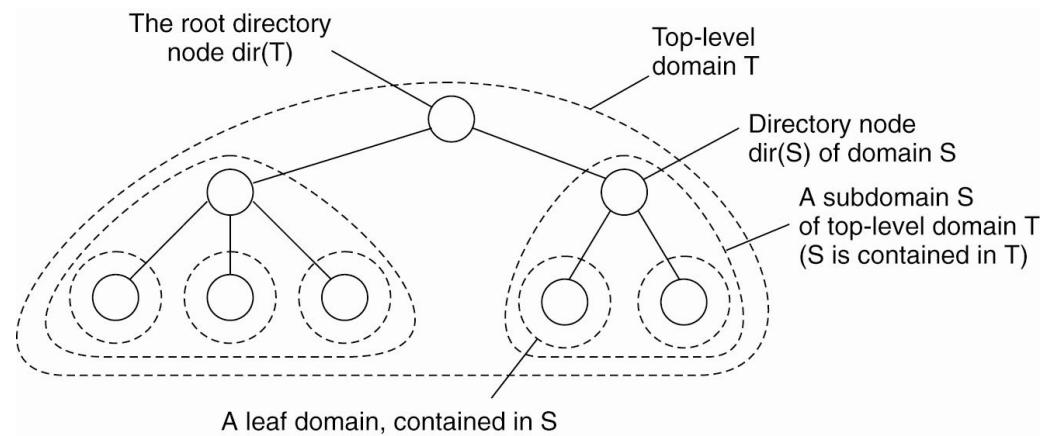
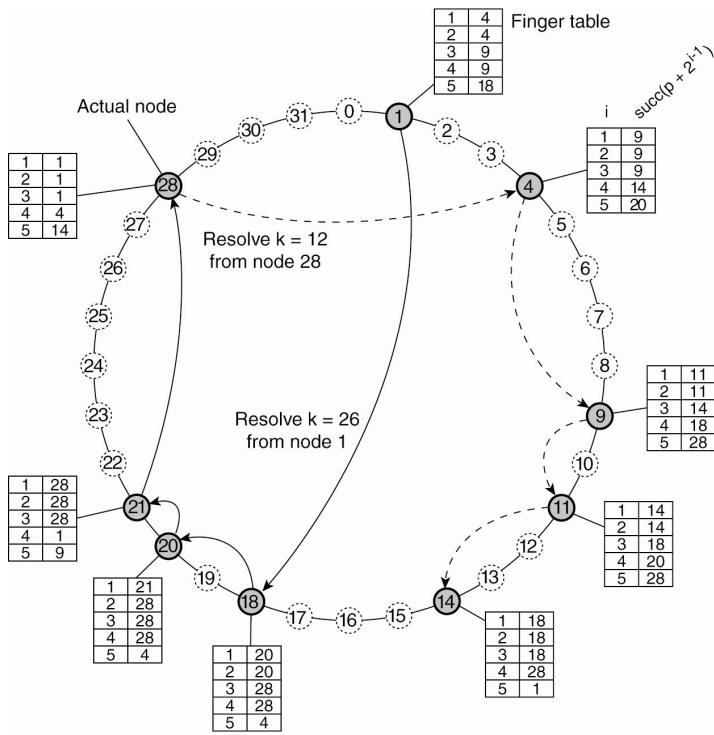


New Topic: Naming

- Names are used to share resources, uniquely identify entities and refer to locations
- Need to map from name to the entity it refers to
 - E.g., Browser access to www.cnn.com
 - Use name resolution
- Differences in naming in distributed and non-distributed systems
 - Distributed systems: naming systems is itself distributed
- How to name mobile entities?



Approaches

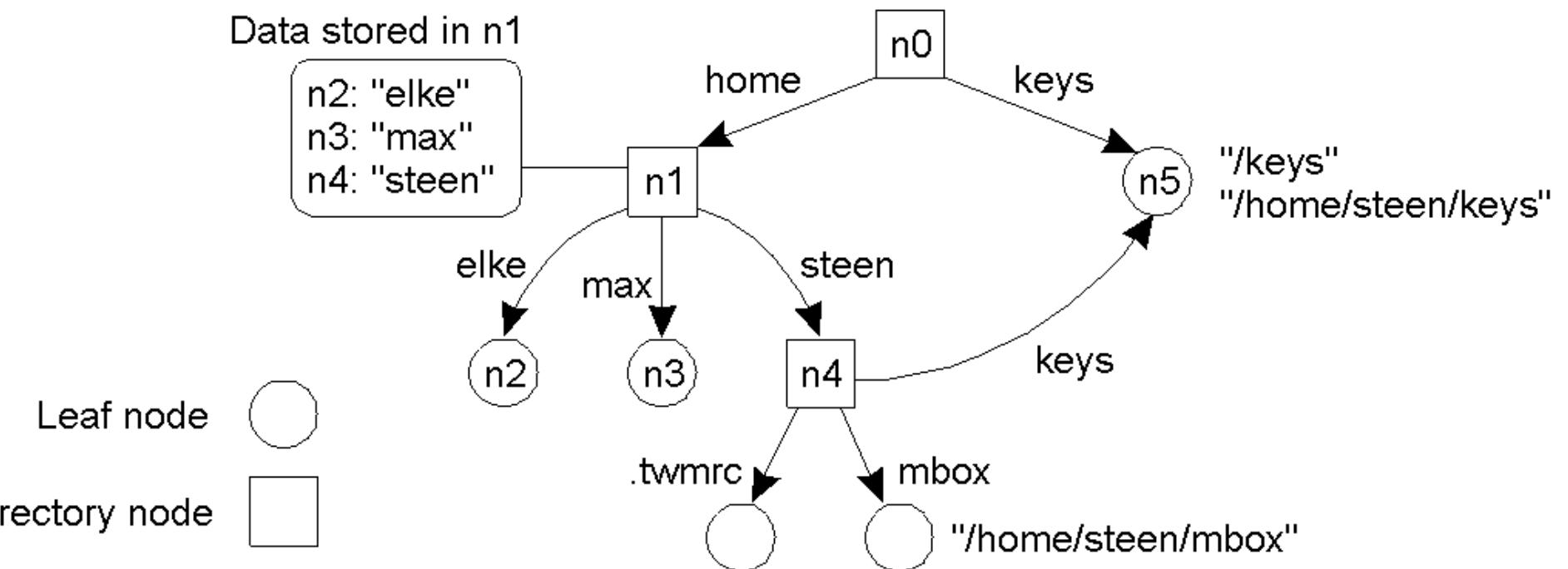


- Hierarchical versus hash-based



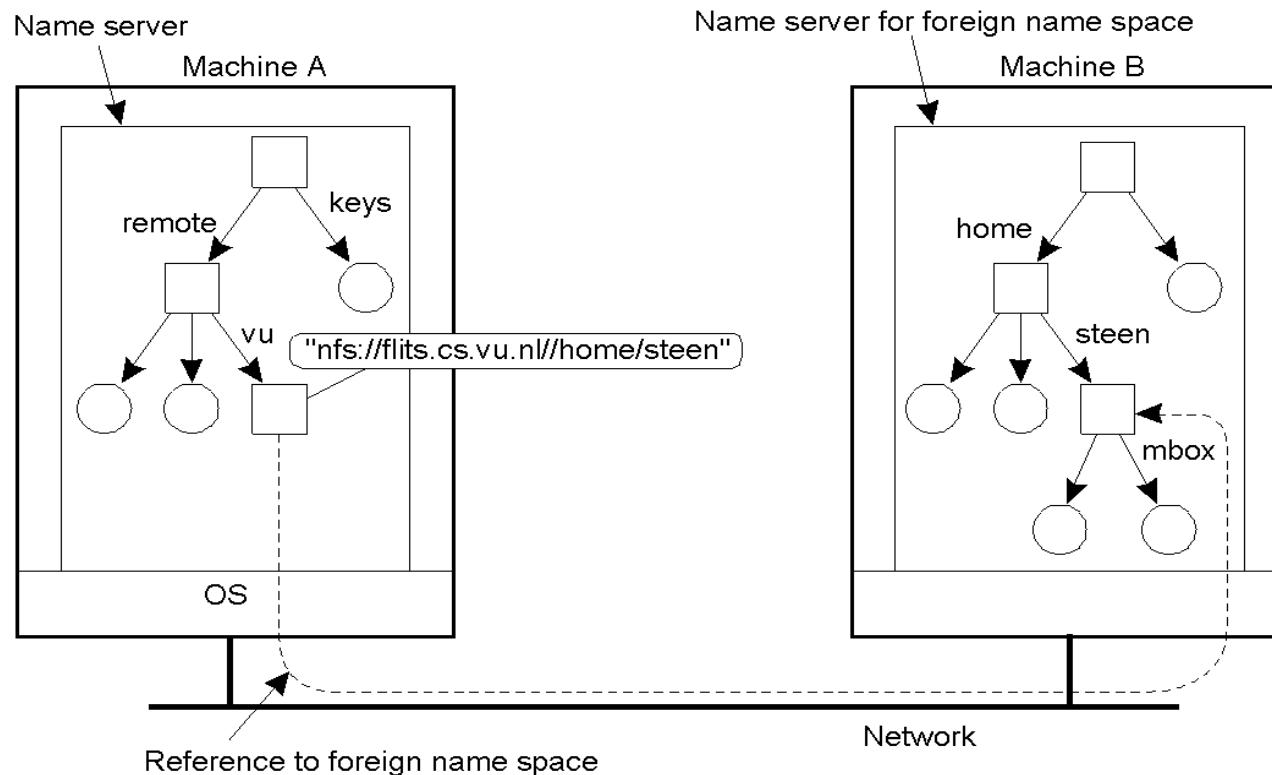
Example: File Names, URLs

- Objects have names; names need to be resolved to map them to the object.
- Hierarchical directory structure (DAG)
 - Each file name is a unique path in the DAG
 - Resolution of `/home/steen/mbox` a traversal of the DAG
- File names are *human-friendly*



Resolving File Names across Machines

- Remote files are accessed using a node name, path name
- NFS mount protocol: map a remote node onto local DAG
 - Remote files are accessed using local names! (*location independence*)
 - OS maintains a mount table with the mappings

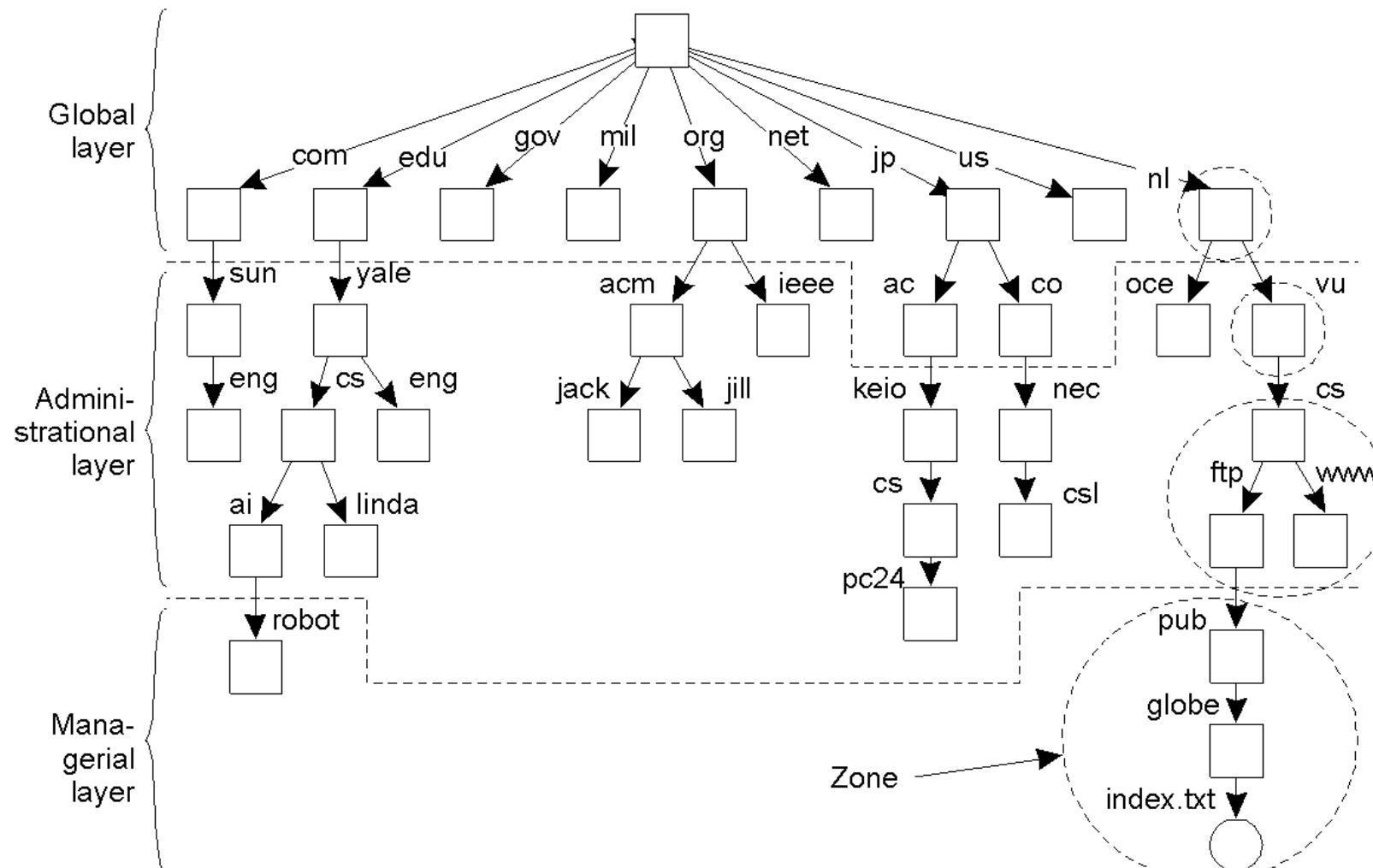


Name Space Distribution

- Naming in large distributed systems
 - System may be global in scope (e.g., Internet, WWW)
- Name space is organized hierarchically
 - Single root node (like naming files)
- Name space is distributed and has three logical layers
 - Global layer: highest level nodes (root and a few children)
 - Represent groups of organizations, rare changes
 - Administrational layer: nodes managed by a single organization
 - Typically one node per department, infrequent changes
 - Managerial layer: actual nodes
 - Frequent changes
 - Zone: part of the name space managed by a separate name server



Name Space Distribution Example



- An example partitioning of the DNS name space, including Internet-accessible files, into three layers.



Name Space Distribution

| Item | Global | Administrational | Managerial |
|---------------------------------|-----------|------------------|--------------|
| Geographical scale of network | Worldwide | Organization | Department |
| Total number of nodes | Few | Many | Vast numbers |
| Responsiveness to lookups | Seconds | Milliseconds | Immediate |
| Update propagation | Lazy | Immediate | Immediate |
| Number of replicas | Many | None or few | None |
| Is client-side caching applied? | Yes | Yes | Sometimes |

- A comparison between name servers for implementing nodes from a large-scale name space partitioned into a global layer, as an administrative layer, and a managerial layer.
- The more stable a layer, the longer are the lookups valid (and can be cached longer)



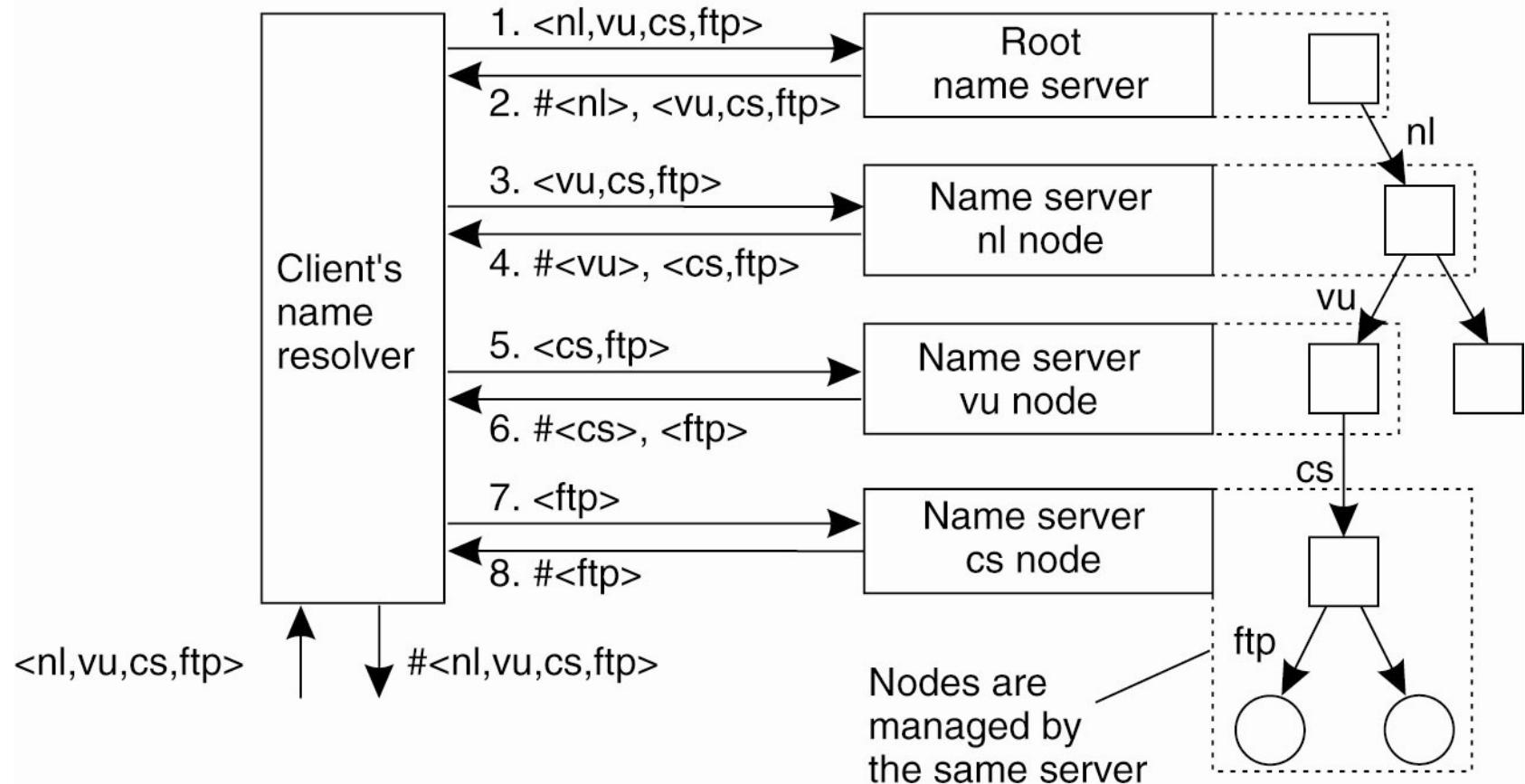
The DNS Name Space

| Type of record | Associated entity | Description |
|----------------|-------------------|---|
| SOA | Zone | Holds information on the represented zone |
| A | Host | Contains an IP address of the host this node represents |
| MX | Domain | Refers to a mail server to handle mail addressed to this node |
| SRV | Domain | Refers to a server handling a specific service |
| NS | Zone | Refers to a name server that implements the represented zone |
| CNAME | Node | Symbolic link with the primary name of the represented node |
| PTR | Host | Contains the canonical name of a host |
| HINFO | Host | Holds information on the host this node represents |
| TXT | Any kind | Contains any entity-specific information considered useful |

- The most important types of resource records forming the contents of nodes in the DNS name space.



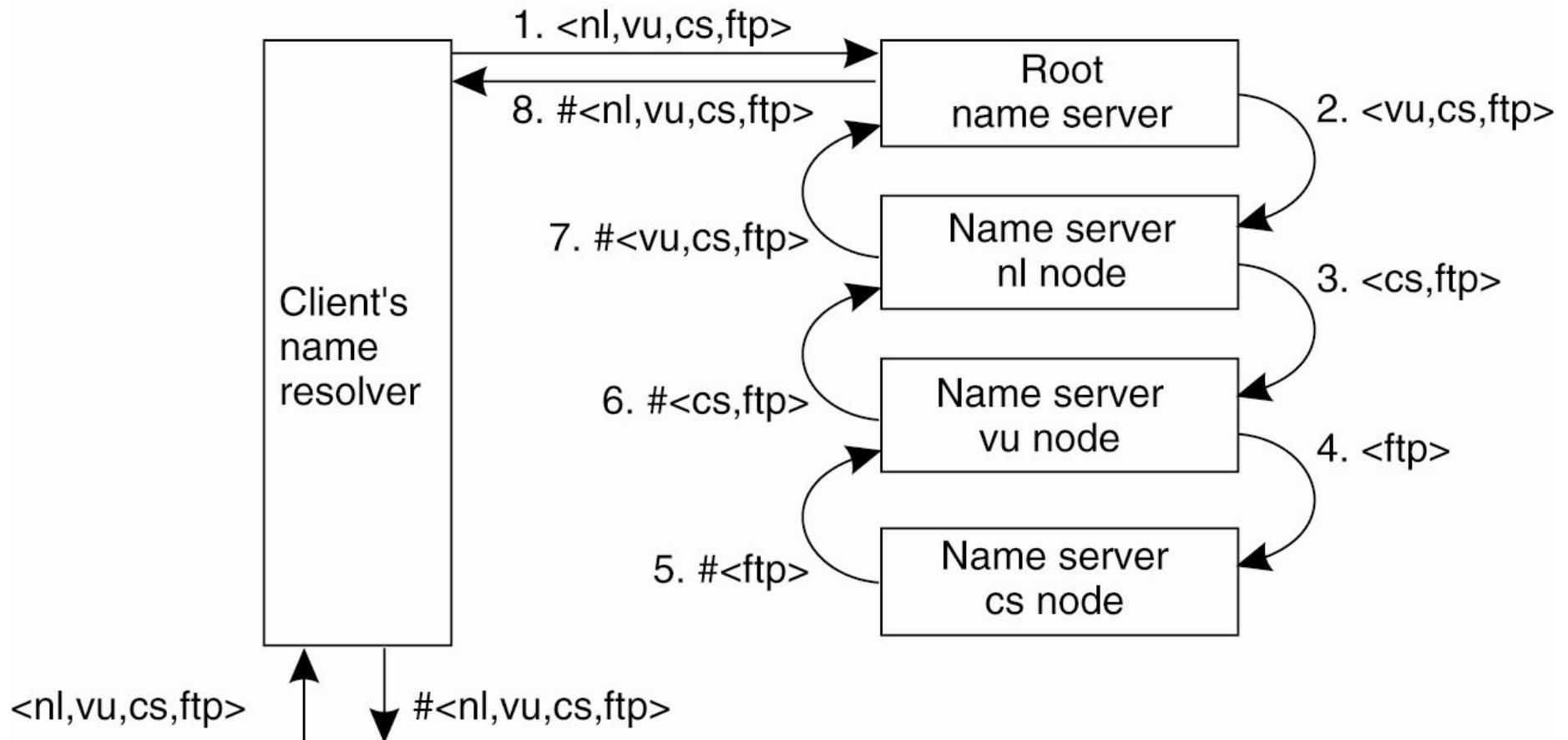
Iterative Resolution



- The principle of iterative name resolution.



Recursive Resolution



- The principle of recursive name resolution.



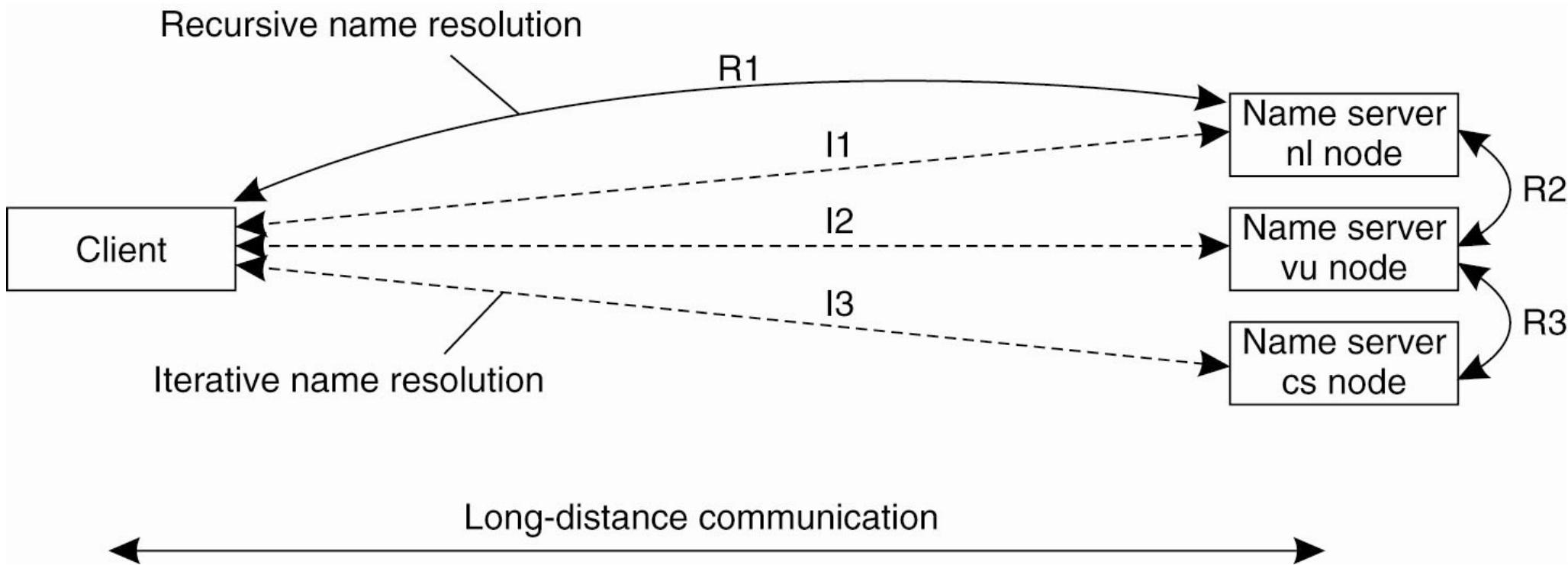
Implementation of Name Resolution

| Server for node | Should resolve | Looks up | Passes to child | Receives and caches | Returns to requester |
|-----------------|----------------|----------|-----------------|-----------------------------------|---|
| cs | <ftp> | #<ftp> | — | — | #<ftp> |
| vu | <cs,ftp> | #<cs> | <ftp> | #<ftp> | #<cs> #<cs, ftp> |
| nl | <vu,cs,ftp> | #<vu> | <cs,ftp> | #<cs> #<cs,ftp> | #<vu> #<vu,cs> #<vu,cs,ftp> |
| root | <nl,vu,cs,ftp> | #<nl> | <vu,cs,ftp> | #<vu> #<vu,cs> #<vu,cs,ftp> | #<nl> #<nl,vu> #<nl,vu,cs> #<nl,vu,cs,ftp> |

- Recursive name resolution of $\langle nl, vu, cs, ftp \rangle$. Name servers cache intermediate results for subsequent lookups.



DNS Resolution



- Comparison between recursive and iterative name resolution with respect to communication costs.



DNS Implementation

- An excerpt from the DNS database for the zone *cs.vu.nl.*

| Name | Record type | Record value |
|-------------------|-------------|---|
| cs.vu.nl | SOA | star (1999121502,7200,3600,2419200,86400) |
| cs.vu.nl | NS | star.cs.vu.nl |
| cs.vu.nl | NS | top.cs.vu.nl |
| cs.vu.nl | NS | solo.cs.vu.nl |
| cs.vu.nl | TXT | "Vrije Universiteit - Math. & Comp. Sc." |
| cs.vu.nl | MX | 1 zephyr.cs.vu.nl |
| cs.vu.nl | MX | 2 tornado.cs.vu.nl |
| cs.vu.nl | MX | 3 star.cs.vu.nl |
| star.cs.vu.nl | HINFO | Sun Unix |
| star.cs.vu.nl | MX | 1 star.cs.vu.nl |
| star.cs.vu.nl | MX | 10 zephyr.cs.vu.nl |
| star.cs.vu.nl | A | 130.37.24.6 |
| star.cs.vu.nl | A | 192.31.231.42 |
| zephyr.cs.vu.nl | HINFO | Sun Unix |
| zephyr.cs.vu.nl | MX | 1 zephyr.cs.vu.nl |
| zephyr.cs.vu.nl | MX | 2 tornado.cs.vu.nl |
| zephyr.cs.vu.nl | A | 192.31.231.66 |
| www.cs.vu.nl | CNAME | soling.cs.vu.nl |
| ftp.cs.vu.nl | CNAME | soling.cs.vu.nl |
| soling.cs.vu.nl | HINFO | Sun Unix |
| soling.cs.vu.nl | MX | 1 soling.cs.vu.nl |
| soling.cs.vu.nl | MX | 10 zephyr.cs.vu.nl |
| soling.cs.vu.nl | A | 130.37.24.11 |
| laser.cs.vu.nl | HINFO | PC MS-DOS |
| laser.cs.vu.nl | A | 130.37.30.32 |
| vucs-das.cs.vu.nl | PTR | 0.26.37.130.in-addr.arpa |
| vucs-das.cs.vu.nl | A | 130.37.26.0 |



X.500 Directory Service

- OSI Standard
- Directory service: special kind of naming service where:
 - Clients can lookup entities based on attributes instead of full name
 - Real-world example: Yellow pages: look for a plumber



LDAP

- Lightweight Directory Access Protocol (LDAP)
 - X.500 too complex for many applications
 - LDAP: Simplified version of X.500
 - Widely used for Internet services
 - Application-level protocol, uses TCP
 - Lookups and updates can use strings instead of OSI encoding
 - Use master servers and replicas servers for performance improvements
 - Example LDAP implementations:
 - Active Directory (Windows 2000)
 - Novell Directory services
 - iPlanet directory services (Netscape)
 - OpenLDAP
 - Typical uses: user profiles, access privileges, network resources



The LDAP Name Space

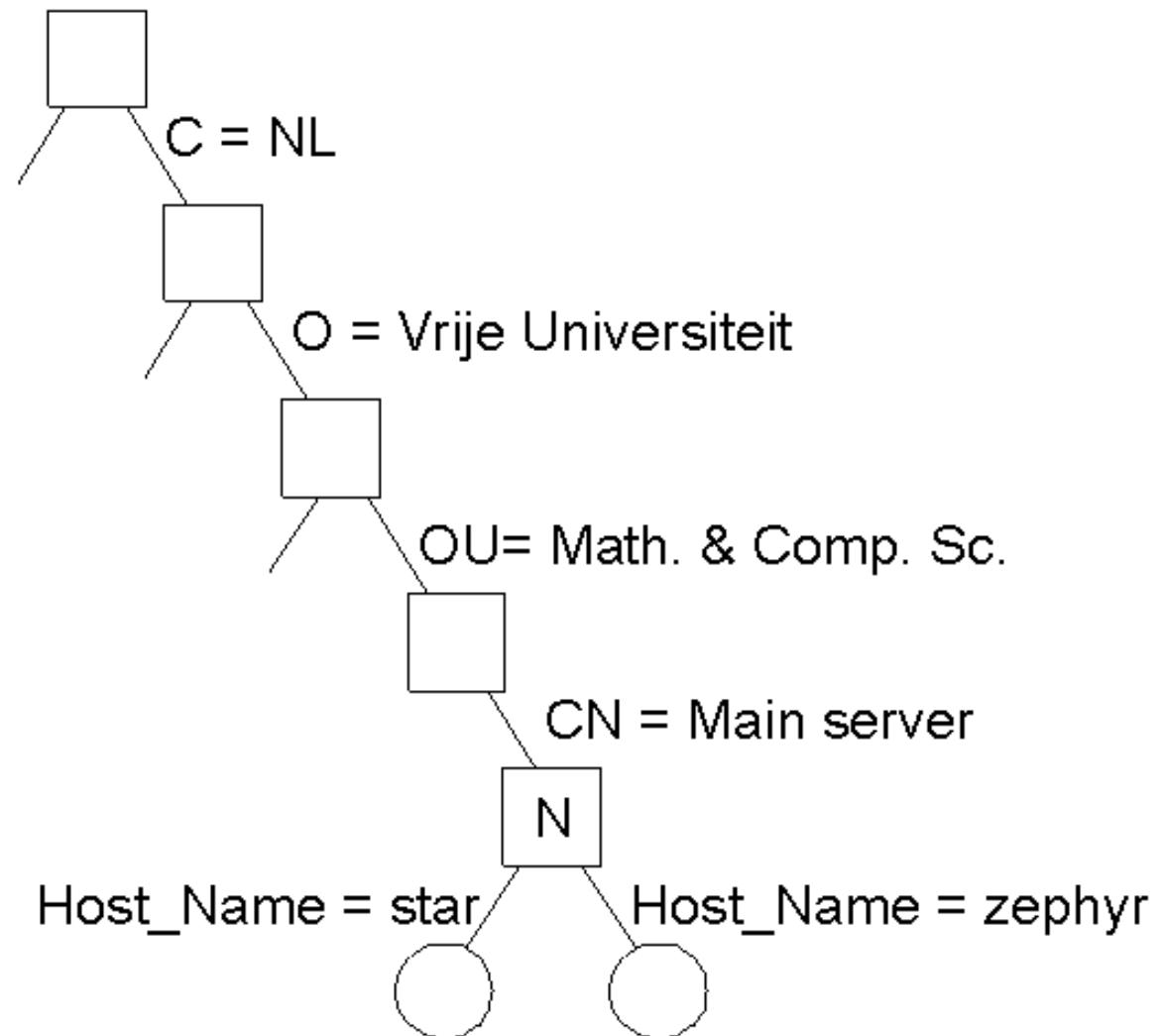
| Attribute | Abbr. | Value |
|--------------------|-------|---------------------------------------|
| Country | C | NL |
| Locality | L | Amsterdam |
| Organization | L | Vrije Universiteit |
| OrganizationalUnit | OU | Math. & Comp. Sc. |
| CommonName | CN | Main server |
| Mail_Servers | -- | 130.37.24.6, 192.31.231,192.31.231.66 |
| FTP_Server | -- | 130.37.21.11 |
| WWW_Server | -- | 130.37.21.11 |

- A simple example of a LDAP directory entry using X.500 naming conventions.

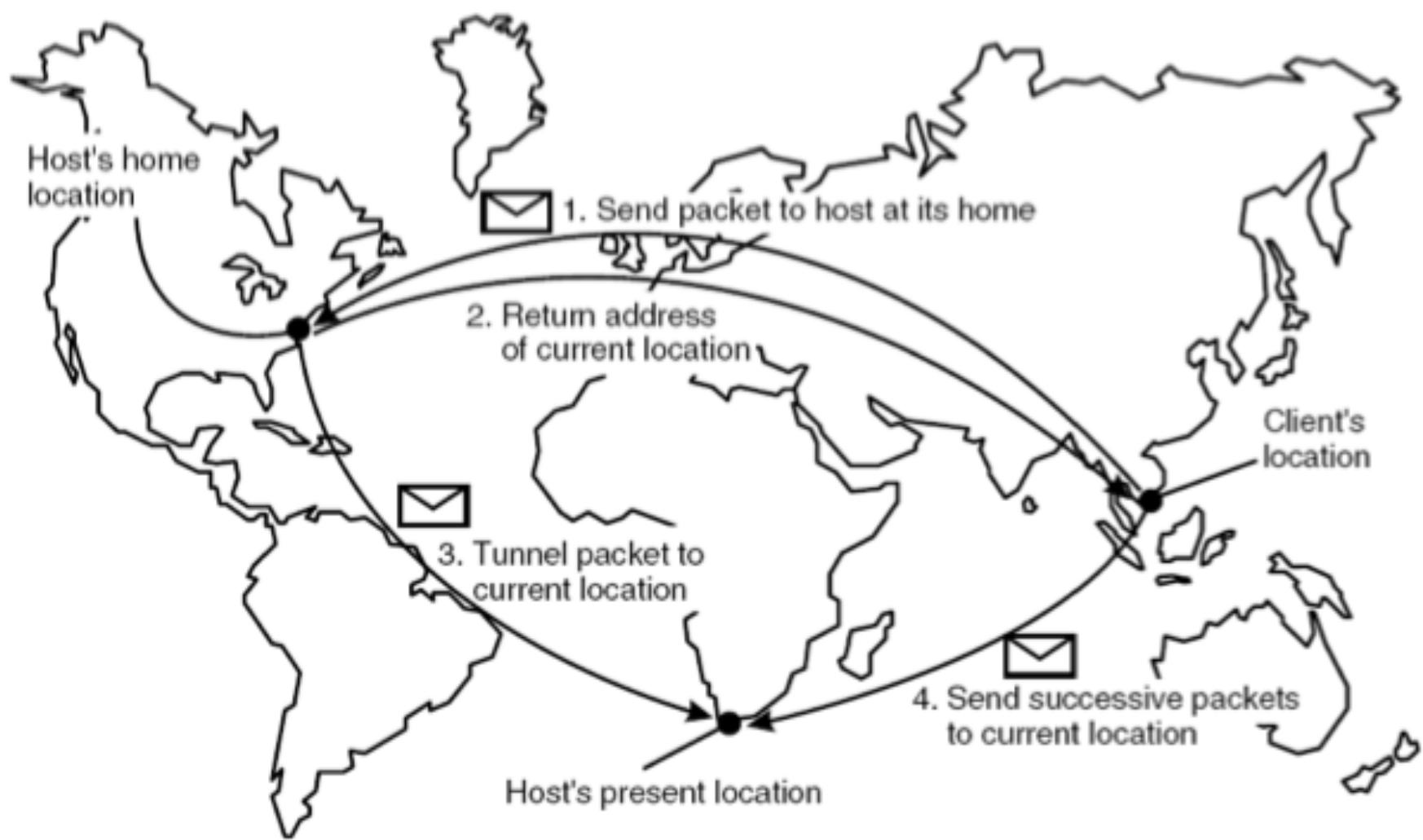


The LDAP Name Space (2)

- Part of the directory information tree.



Home-Based Approaches



The principle of Mobile IP.

Computer Science

CS677: Distributed OS

Lecture 10, page 18