Epidemic Protocols

• Used in Bayou system from Xerox PARC
• Bayou: weakly connected replicas
  – Useful in mobile computing (mobile laptops)
  – Useful in wide area distributed databases (weak connectivity)
• Based on theory of epidemics (spreading infectious diseases)
  – Upon an update, try to “infect” other replicas as quickly as possible
  – Pair-wise exchange of updates (like pair-wise spreading of a disease)
  – Terminology:
    • Infective store: store with an update it is willing to spread
    • Susceptible store: store that is not yet updated
• Many algorithms possible to spread updates

Spreading an Epidemic

• Anti-entropy
  – Server $P$ picks a server $Q$ at random and exchanges updates
  – Three possibilities: only push, only pull, both push and pull
  – Claim: A pure push-based approach does not help spread updates quickly (Why?)
    • Pull or initial push with pull work better
• Rumor mongering (aka gossiping)
  – Upon receiving an update, $P$ tries to push to $Q$
  – If $Q$ already received the update, stop spreading with prob $1/k$
  – Analogous to “hot” gossip items $\Rightarrow$ stop spreading if “cold”
  – Does not guarantee that all replicas receive updates
    • Chances of staying susceptible: $s = e^{-(k+1)(1-s)}$
Removing Data

- Deletion of data items is hard in epidemic protocols
- Example: server deletes data item $x$
  - No state information is preserved
    - Can’t distinguish between a deleted copy and no copy!
- Solution: death certificates
  - Treat deletes as updates and spread a death certificate
    - Mark copy as deleted but don’t delete
    - Need an eventual clean up
      - Clean up dormant death certificates

Implementation Issues

- Two techniques to implement consistency models
  - Primary-based protocols
    - Assume a primary replica for each data item
    - Primary responsible for coordinating all writes
  - Replicated write protocols
    - No primary is assumed for a data item
    - Writes can take place at any replica
Remote-Write Protocols

- Traditionally used in client-server systems

Remote-Write Protocols (2)

- Primary-backup protocol
  - Allow local reads, sent writes to primary
  - Block on write until all replicas are notified
  - Implements sequential consistency
Local-Write Protocols (1)

- Primary-based local-write protocol in which a single copy is migrated between processes.
  - Limitation: need to track the primary for each data item

1. Read or write request
2. Forward request to current server for x
3. Move item x to client's server
4. Return result of operation on client's server

Local-Write Protocols (2)

- Primary-backup protocol in which the primary migrates to the process wanting to perform an update

W1. Write request
W2. Move item x to new primary
W3. Acknowledge write completed
W4. Tell backups to update
W5. Acknowledge update

R1. Read request
R2. Response to read
Replicated-write Protocols

- Relax the assumption of one primary
  - No primary, any replica is allowed to update
  - Consistency is more complex to achieve

- Quorum-based protocols
  - Use voting to request/acquire permissions from replicas
  - Consider a file replicated on $N$ servers
    - $N_R + N_W > N$, $N_W > N/2$
    - Update: contact at least $(N/2+1)$ servers and get them to agree to do update (associate version number with file)
    - Read: contact majority of servers and obtain version number
    - If majority of servers agree on a version number, read

Gifford’s Quorum-Based Protocol

- Three examples of the voting algorithm:
  a) A correct choice of read and write set
  b) A choice that may lead to write-write conflicts
  c) A correct choice, known as ROWA (read one, write all)
Replica Management

• Replica server placement
  – Web: geophysically skewed request patterns
  – Where to place a proxy?
    • K-clusters algorithm

• Permanent replicas versus temporary
  – Mirroring: all replicas mirror the same content
  – Proxy server: on demand replication

• Server-initiated versus client-initiated

Content Distribution

• Will come back to this in Chap 12

• CDN: network of proxy servers
• Caching:
  – update versus invalidate
  – Push versus pull-based approaches
  – Stateful versus stateless
• Web caching: what semantics to provide?
Final Thoughts

- Replication and caching improve performance in distributed systems
- Consistency of replicated data is crucial
- Many consistency semantics (models) possible
  - Need to pick appropriate model depending on the application
  - Example: web caching: weak consistency is OK since humans are tolerant to stale information (can reload browser)
  - Implementation overheads and complexity grows if stronger guarantees are desired

Fault Tolerance

- Single machine systems
  - Failures are all or nothing
    - OS crash, disk failures
- Distributed systems: multiple independent nodes
  - Partial failures are also possible (some nodes fail)
- Question: Can we automatically recover from partial failures?
  - Important issue since probability of failure grows with number of independent components (nodes) in the systems
  - \[ \text{Prob(failure)} = \text{Prob(Any one component fails)} = 1 - \text{P(no failure)} \]
A Perspective

- Computing systems are not very reliable
  - OS crashes frequently (Windows), buggy software, unreliable hardware, software/hardware incompatibilities
  - Until recently: computer users were “tech savvy”
    - Could depend on users to reboot, troubleshoot problems
  - Growing popularity of Internet/World Wide Web
    - “Novice” users
    - Need to build more reliable/dependable systems
  - Example: what is your TV (or car) broke down every day?
    - Users don’t want to “restart” TV or fix it (by opening it up)
- Need to make computing systems more reliable
  - Important for online banking, e-commerce, online trading, webmail…

Basic Concepts

- Need to build *dependable* systems
- Requirements for dependable systems
  - Availability: system should be available for use at any given time
    - 99.999 % availability (five 9s) ⇒ very small down times
  - Reliability: system should run continuously without failure
  - Safety: temporary failures should not result in a catastrophic
    - Example: computing systems controlling an airplane, nuclear reactor
  - Maintainability: a failed system should be easy to repair
Basic Concepts (contd)

- Fault tolerance: system should provide services despite faults
  - Transient faults
  - Intermittent faults
  - Permanent faults

## Failure Models

<table>
<thead>
<tr>
<th>Type of failure</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Crash failure</td>
<td>A server halts, but is working correctly until it halts</td>
</tr>
<tr>
<td>Omission failure</td>
<td>A server fails to respond to incoming requests</td>
</tr>
<tr>
<td></td>
<td>Receive omission</td>
</tr>
<tr>
<td></td>
<td>A server fails to receive incoming messages</td>
</tr>
<tr>
<td></td>
<td>Send omission</td>
</tr>
<tr>
<td></td>
<td>A server fails to send messages</td>
</tr>
<tr>
<td>Timing failure</td>
<td>A server’s response lies outside the specified time interval</td>
</tr>
<tr>
<td>Response failure</td>
<td>The server's response is incorrect</td>
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<tr>
<td></td>
<td>Value failure</td>
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<tr>
<td></td>
<td>The value of the response is wrong</td>
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<tr>
<td></td>
<td>State transition failure</td>
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<tr>
<td></td>
<td>The server deviates from the correct flow of control</td>
</tr>
<tr>
<td>Arbitrary failure</td>
<td>A server may produce arbitrary responses at arbitrary times</td>
</tr>
</tbody>
</table>

- Different types of failures.
Failure Masking by Redundancy

- Triple modular redundancy.