Today: Synchronization for Readers/Writers Problem

• An object is shared among may threads, each belonging to one of two classes:
  – Readers: read data, never modify it
  – Writers: read data and modify it
• Using a single lock on the data object is overly restrictive
  \(\Rightarrow\) Want many readers reading the object at once
  – Allow only one writer at any point
  – How do we control access to the object to permit this protocol?
• Correctness criteria:
  – Each read or write of the shared data must happen within a critical section.
  – Guarantee mutual exclusion for writers.
  – Allow multiple readers to execute in the critical section at once.

Readers/Writers Problem

class ReadWrite {
    public:
        void Read();
        void Write();
    private:
        int readers; // counts readers
        Semaphore mutex; // controls access to readers
        Semaphore wrt; // controls entry to first
    } // writer or reader
ReadWrite::ReadWrite {
    readers = 0;
    mutex->value = 1;
    wrt->value = 1;
}
Readers/Writers Problem

```
ReadWrite::Write()
{
    wrt.wait();    // any writers or readers?
    <perform write>
    wrt.Signal();  // enable others
}

ReadWrite::Read()
{
    mutex.Wait();   // ensure mutual exclusion
    readers += 1;   // another reader
    if (readers == 1)
        wrt->Wait(); // block writers
    mutex.Signal();
    <perform read>
    mutex.Wait();   // ensure mutual exclusion
    readers -= 1;   // reader done
    if (readers == 0)
        wrt.Signal(); // enable writers
    mutex.Signal();
}
```

Readers/Writers: Scenario 1

<table>
<thead>
<tr>
<th>R1:</th>
<th>R2:</th>
<th>W1:</th>
</tr>
</thead>
<tbody>
<tr>
<td>Read ()</td>
<td>Read ()</td>
<td>Write ()</td>
</tr>
</tbody>
</table>

Readers/Writers: Scenario 2

R1:                R2:               W1:
    Read ()            Read ()        Write ()
    Read ()            Read ()

Reader/Writers: Scenario 3

R1:                R2:               W1:
    Read ()            Read ()        Write ()
    Read ()
Readers/Writers Solution: Discussion

- Implementation notes:
  1. The first reader blocks if there is a writer; any other readers who try to enter block on mutex.
  2. The last reader to exit signals a waiting writer.
  3. When a writer exits, if there is both a reader and writer waiting, which goes next depends on the scheduler.
  4. If a writer exits and a reader goes next, then all readers that are waiting will fall through (at least one is waiting on wrt and zero or more can be waiting on mutex).
  5. Does this solution guarantee all threads will make progress?

- Alternative desirable semantics:
  - Let a writer enter its critical section as soon as possible.

```
ReadWrite::Write()
{
    write_mutex.Wait(); // ensure mutual exclusion
    writers += 1; // another pending writer
    if (writers == 1) // block readers
        read_block.Wait();
    write_mutex.Signal();
    write_block.Wait(); // ensure mutual exclusion
    <perform write>
    write_block.Signal();
    write_mutex.Wait(); // ensure mutual exclusion
    writers -= 1; // writer done
    if (writers == 0) // enable readers
        read_block.Signal();
    write_mutex.Signal(); } 
```
Readers/Writers Solution Favoring Writers

```c
ReadWrite::Read()

  write_pending->Wait();  // ensures at most one reader will go
  // before a pending write
read_block->Wait();  // write in progress; wait
read_mutex->Wait();   // ensure mutual exclusion
readers += 1;         // another reader
if (readers == 1)     // synchronize with writers
  write_block->Wait();
read_mutex->Signal();
read_block->Signal();
write_pending->Signal();
<perform read>
read_mutex->Wait();   // ensure mutual exclusion
readers -= 1;         // reader done
if (readers == 0)     // enable writers
  write_block->Signal();
read_mutex->Signal(); }
```

Readers/Writers: Scenario 4

<table>
<thead>
<tr>
<th>R1:</th>
<th>R2:</th>
<th>W1:</th>
<th>W2:</th>
</tr>
</thead>
<tbody>
<tr>
<td>Read ()</td>
<td>Read ()</td>
<td>Write ()</td>
<td>Write ()</td>
</tr>
</tbody>
</table>
Readers/Writers: Scenario 5

R1:                R2:               W1:              W2:
Read ()          Read ()      Write ()        Write ()

Reader/Writers: Scenario 6

R1:                R2:               W1:              W2:
Read ()          Read ()      Write ()        Write ()
Readers/Writers using Monitors (Java)

class ReaderWriter {
    private int numReaders = 0;
    private int numWriters = 0;

    private synchronized void prepareToRead () {
        while ( numWriters > 0 ) wait ();
        numReaders++;
    }

    private synchronized void doneReading () {
        numReaders--;
        if ( numReaders == 0 ) notify ();
    }

    public ... someReadMethod () {
        // reads NOT synchronized: multiple readers
        prepareToRead ();
        <do the reading>
        doneReading ();
    }
}

private synchronized void
    doneReading () {
        numReaders--;
        if ( numReaders == 0 ) notify ();
    }

public ... someReadMethod () {
    // reads NOT synchronized: multiple readers
    prepareToRead ();
    <do the reading>
    doneReading ();
}

private synchronized void
    doneWriting () {
        numWriters--;
        notify ();
    }

private synchronized void someWriteMethod (...) {
    // synchronized => only one writer
    prepareToWrite ();
    <do the writing>
    doneWriting ();
}

private synchronized void
    doneWriting () {
        numWriters--;
        notify ();
    }

public synchronized void someWriteMethod (...) {
    // synchronized => only one writer
    prepareToWrite ();
    <do the writing>
    doneWriting ();
}
Read/write Locks

- pthreads and Java support read/write locks
  - A thread can acquire a read lock or a write lock
- Multiple threads can hold the same read lock concurrently
- Only one thread can hold a write lock
- Java: ReadWriteLock class
  - readLock()
  - writeLock()
- pthread routines:
  - pthread_rwlock_init()
  - pthread_rwlock_rdlock()
  - pthread_rwlock_wrlock()
  - pthread_rwlock_unlock()

Dining Philosophers

- It’s lunch time in the philosophy dept
- Five philosophers, each either eats or thinks
- Share a circular table with five chopsticks
- Thinking: do nothing
- Eating => need two chopsticks, try to pick up two closest chopsticks
  - Block if neighbor has already picked up a chopstick
- After eating, put down both chopsticks and go back to thinking
Dining Philosophers v1

Semaphore chopsticks[5];

do{
    wait(chopstick[i]);  // left chopstick
    wait(chopstick[(i+1)%5 ]); // right chopstick
    // eat
    signal(chopstick[i]);  // left chopstick
    signal(chopstick[(i+1)%5 ]); // right chopstick
    // think
} while(TRUE);

Dining Philosophers v2 (monitors)

monitor DP
{
    enum { THINKING; HUNGRY, EATING) state [5] ;
    condition self [5];

    void synchronized pickup (int i) {
        state[i] = HUNGRY;
        test(i);
        if (state[i] != EATING)
            self[i].wait;
    }

    void synchronized putdown (int i) {
        state[i] = THINKING;
        //test left and right neighbors
        test((i + 4) % 5);
        test((i + 1) % 5);
    }

    void test (int i) {
        if ( (state[(i + 4) % 5] != EATING) &&
            (state[i] == HUNGRY) &&
            (state[(i + 1) % 5] != EATING) ) {
            state[i] = EATING ;
            self[i].signal () ;
        }
    }

    initialization_code() {
        for (int i = 0; i < 5; i++)
            state[i] = THINKING;
    }
}
Dining Philosophers (semaphores)

```c
#define N 5 /* number of philosophers */
#define LEFT (i+N-1)%N /* number of i’s left neighbor */
#define RIGHT (i+1)%N /* number of i’s right neighbor */
#define THINKING 0 /* philosopher is thinking */
#define HUNGRY 1 /* philosopher is trying to get forks */
#define EATING 2 /* philosopher is eating */
typedef int semaphore;
int state[N];
semaphore mutex = 1;
semaphore s[N];

void philosopher(int i) /* i: philosopher number, from 0 to N−1 */
{
    while (TRUE) { /* repeat forever */
        think(); /* philosopher is thinking */
        take_forks(i); /* acquire two forks or block */
        eat(); /* yum-yum, spaghetti */
        put_forks(i); /* put both forks back on table */
    }
}
```

Dining Philosophers (contd)

```c
void take_forks(int i) /* i: philosopher number, from 0 to N−1 */
{
    down(&mutex); /* enter critical region */
    state[i] = HUNGRY; /* record fact that philosopher i is hungry */
    test(i); /* try to acquire 2 forks */
    up(&mutex); /* exit critical region */
    down(&s[i]); /* block if forks were not acquired */
}

void put_forks(i) /* i: philosopher number, from 0 to N−1 */
{
    down(&mutex); /* enter critical region */
    state[i] = THINKING; /* philosopher has finished eating */
    test(LEFT); /* see if left neighbor can now eat */
    test(RIGHT); /* see if right neighbor can now eat */
    up(&mutex); /* exit critical region */
}

void test(i) /* i: philosopher number, from 0 to N−1 */
{
    if (state[i] == HUNGRY && state[LEFT] == EATING && state[RIGHT] != EATING) {
        state[i] = EATING;
        up(&s[i]);
    }
}
```
Real-world Examples

• Producer-consumer
  – Audio-Video player: network and display threads; shared buffer
  – Web servers: master thread and slave thread

• Reader-writer
  – Banking system: read account balances versus update

• Dining Philosophers
  – Cooperating processes that need to share limited resources
    • Set of processes that need to lock multiple resources
      – Disk and tape (backup),
    • Travel reservation: hotel, airline, car rental databases

Summary

• Readers/writers problem:
  – Allow multiple readers to concurrently access a data
  – Allow only one writer at a time

• Two possible solutions using semaphores
  – Favor readers
  – Favor writers

• Starvation is possible in either case!
• Dining philosophers: mutually exclusive access to multiple resources